

Declarative Merging of and Reasoning about Decision Diagrams

Thomas Eiter Thomas Krennwallner Christoph Redl

{eiter,tkren,redl}@kr.tuwien.ac.at



TECHNISCHE
UNIVERSITÄT
WIEN
Vienna University of Technology



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Outline

- 1 Motivation
- 2 Preliminaries: MELD
- 3 Merging of Decision Diagrams
- 4 Reasoning about Decision Diagrams
- 5 Application: DNA Classification
- 6 Conclusion

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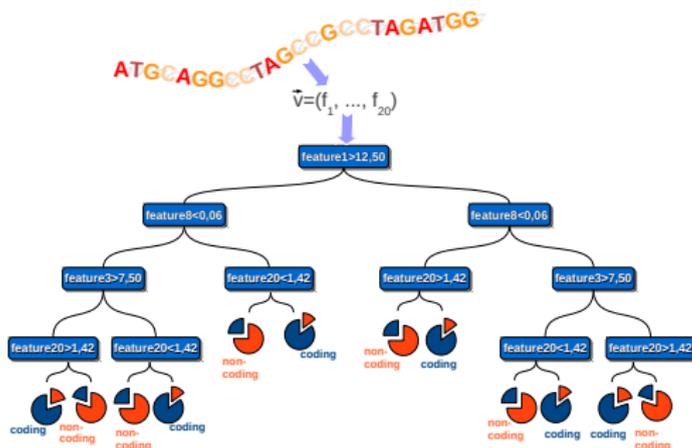
Motivation

Decision Diagrams

- Important means for decision making
- Intuitively understandable
- Not only for knowledge engineers

Examples

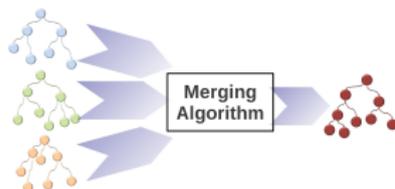
- Severity ratings (e.g. TNM system)
- Diagnosis of personality disorders
- DNA classification



Multiple Diagrams

Reasons

- Different opinions
- Randomized machine-learning algorithms
- Statistical impreciseness



Question: How to combine them?

Multiple Diagram Integration

The DDM System

- Integration process declaratively described
- Ingredients:
 - 1 input decision diagrams
 - 2 merging algorithms
(predefined or user-defined)
- Focus:
 - **process formalization**
 - experimenting with different (combinations of) merging algorithms
 - declarative reasoning for controlling the merging process
- We do **not** focus:
 - concrete merging strategies
 - accuracy improvement

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MELD

Task

- Collection of knowledge bases: $KB = KB_1, \dots, KB_n$
- Associated collections of belief sets: $BS(KB_1), \dots, BS(KB_n) \in \mathbb{B}_\Sigma$
- Goal: **Integrate** them into a single set of belief sets

Method: Merging Operators

$$\circ^{n,m} : \underbrace{(2^{\mathbb{B}_\Sigma})^n}_{\text{collections of belief sets}} \times \underbrace{\mathcal{A}_1 \times \dots \times \mathcal{A}_m}_{\text{operator arguments}} \rightarrow 2^{\mathbb{B}_\Sigma}$$

Example

Operator definition:

$$\circ_{\cup}^{2,0}(\mathcal{B}_1, \mathcal{B}_2) = \{B_1 \cup B_2 \mid B_1 \in \mathcal{B}_1, B_2 \in \mathcal{B}_2, \nexists A : \{A, \neg A\} \subseteq (B_1 \cup B_2)\},$$

Application:

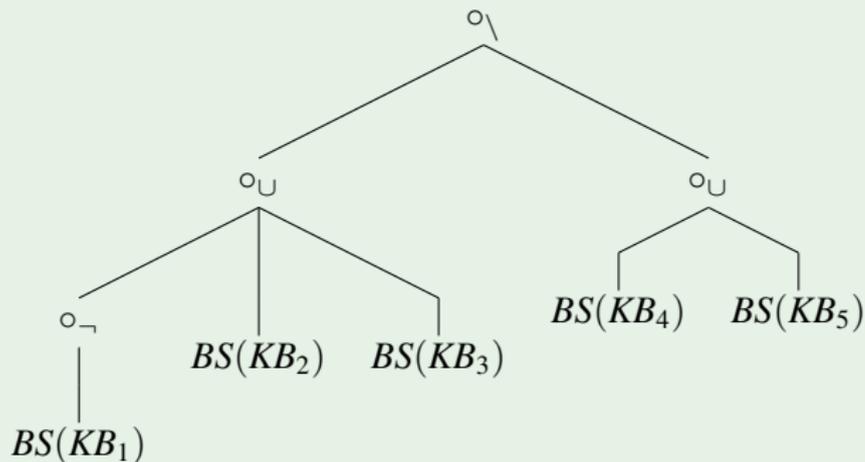
- $\mathcal{B}_1 = \{\{a, b, c\}, \{\neg a, c\}\}, \mathcal{B}_2 = \{\{\neg a, d\}, \{c, d\}\}$
- $\circ_{\cup}^{2,0}(\mathcal{B}_1, \mathcal{B}_2) = \{\{a, b, c, d\}, \{\neg a, c, d\}\}$

MELD

Merging Plan

- Hierarchical arrangement of merging operators

Example



MELD

Merging Tasks

- User provides
 - belief bases with associated collections of belief sets
 - merging plan
 - optional: user-defined merging operators
- MELD: automated evaluation

Advantages

- Reuse of operators
- Quick restructuring of merging plan

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Decision Diagrams

Definition (Decision Diagram)

A **decision diagram** over \mathcal{D} and \mathcal{C} is a labelled rooted directed acyclic graph

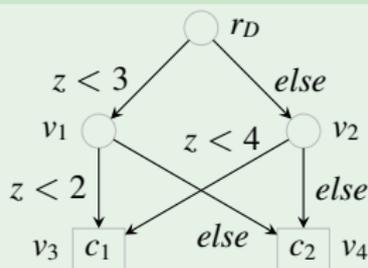
$$D = \langle V, E, \ell_C, \ell_E \rangle$$

- V ... nonempty set of nodes with unique root node $r_D \in V$
- $E \subseteq V \times V$... set of directed edges
- $\ell_C : V \rightarrow \mathcal{C}$... partial function assigning a class to all leafs
- $\ell_E : E \rightarrow \mathcal{Q}$... assign queries $Q(z) : \mathcal{D} \rightarrow \{true, false\}$ to edges
Query language: $O_1 \circ O_2$ with operands O_1, O_2 and $\circ \in \{<, \leq, =, \neq, \geq, >\}$ or "else"

Example

$$\mathcal{D} = \{1, 2, 3, 4, 5\}$$

$$\mathcal{C} = \{c_1, c_2\}$$



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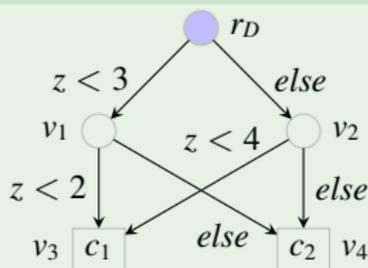
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Classify: 4



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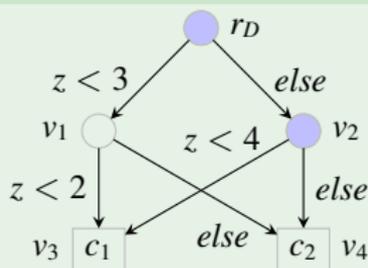
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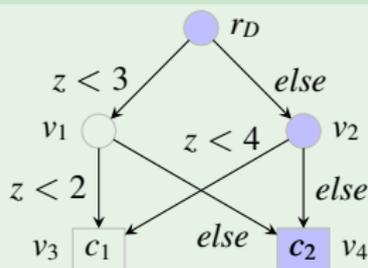
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Classify: $4 \Rightarrow c_2$



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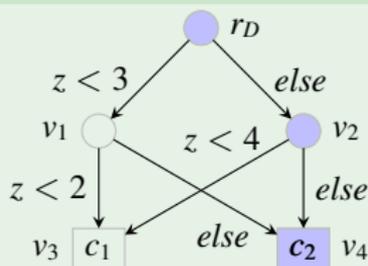
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$$\mathcal{C} = \{c_1, c_2\}$$

Classify: $4 \Rightarrow c_2$



Note: \mathcal{D} may consist of composed objects, e.g. $Q(z) = z.TSH > 4.5mU/l$

Decision Diagram Merging

Instantiation of MELD

- How to use MELD for decision diagram merging?

Decision Diagram Merging

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- How to use MELD for decision diagram merging?
 - 1 Encode decision diagrams as belief sets
 - 2 Merging by special operators

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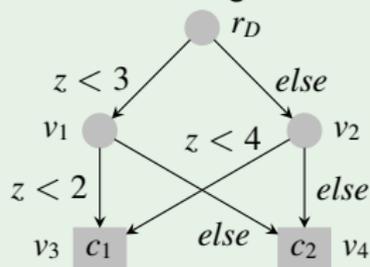
1. Encoding

- Define nodes
 $root(n)$, $inner(n)$, $leaf(n, l)$
- Arcs between nodes, labelled with conditions
 $cond(n_1, n_2, o_1, c, o_2)$, $else(n_1, n_2)$

1. Encoding of Decision Diagrams

Example

Decision Diagram D :

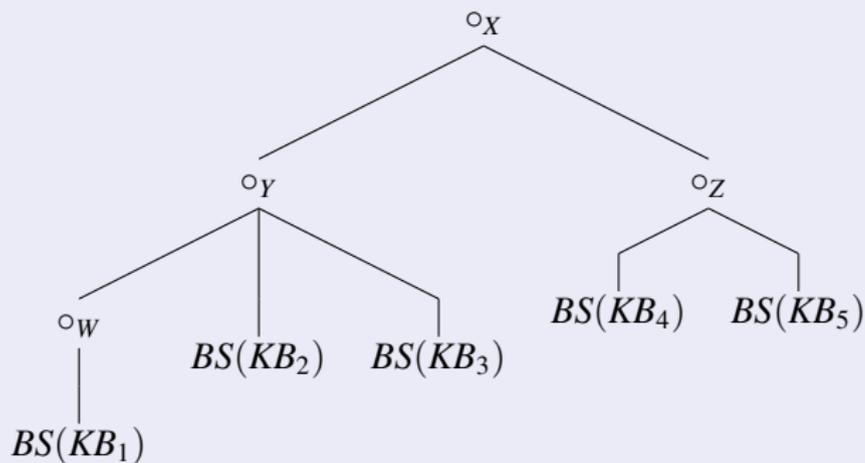


$$\begin{aligned}
 E(D) = \{ & \text{root}(r_D); \text{inner}(r_D); \text{inner}(v_1); \text{inner}(v_2); \\
 & \text{leaf}(v_3, c_1); \text{leaf}(v_4, c_2); \\
 & \text{cond}(r_D, v_1, z, <, 3); \text{else}(r_D, v_2); \\
 & \text{cond}(v_1, v_3, z, <, 2); \text{else}(v_1, v_4); \\
 & \text{cond}(v_2, v_3, z, <, 4); \text{else}(v_2, v_4) \}
 \end{aligned}$$

2. Merging of Decision Diagrams

Merging

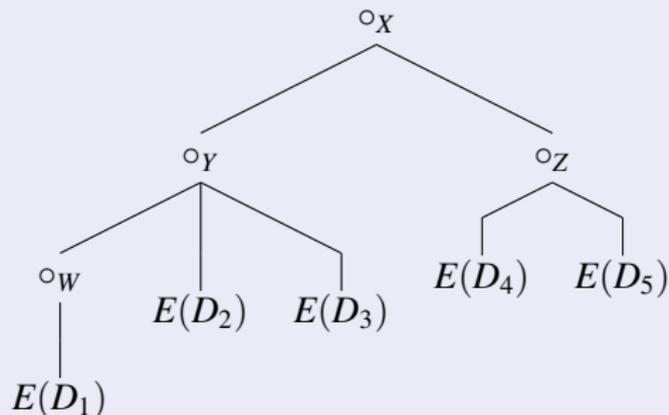
Belief sets = encoded diagrams



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Merging

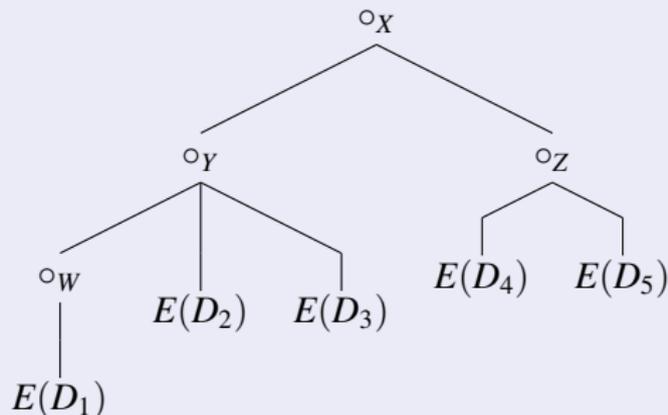
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2. Merging of Decision Diagrams

Merging

Belief sets = encoded diagrams



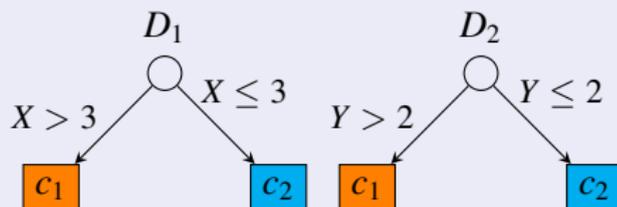
Special merging operators $\circ_W, \circ_X, \circ_Y, \circ_Z$ required!

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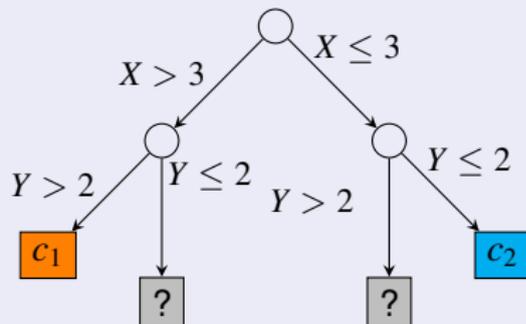
Some Examples of Predefined Operators

■ User Preferences

Give some class label preference over another



$\circ_{pref}(D_1, D_2, c_2 > c_1)$

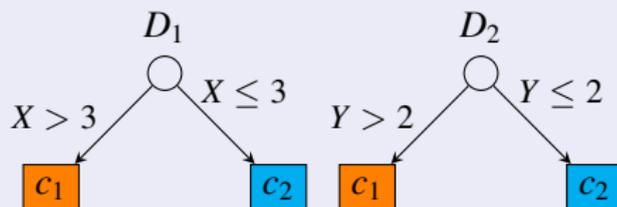


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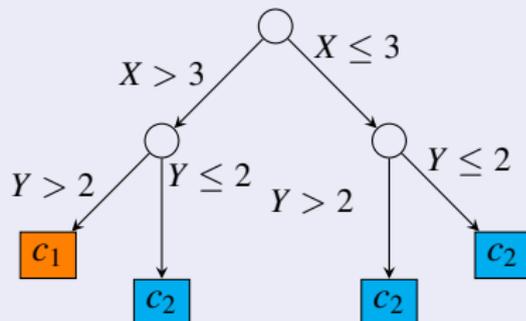
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2. Merging of Decision Diagrams

Some Examples of Predefined Operators

- **User Preferences**
Give some class label preference over another
- **Majority Voting**
Majority of input diagrams decides upon an element's class
- **Simplification**
Decrease redundancy
- **MORGAN merging strategy**
see later
- ...

Note: Operators may produce multiple results!

Example: Majority voting for classes with equal number of votes

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Reasoning about Decision Diagrams

Goal

- Compute **diagram properties**
e.g. height, variable occurrences, redundancy
- Properties may **control** the **merging process** by **filtering**

Reasoning about Decision Diagrams

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Realization

- Special unary operator

$$\circ_{asp}(\Delta, P),$$

Δ ... set of decision diagrams

P ... ASP program

- $P' := P \cup \bigcup_{D \in \Delta} \hat{E}(D)$

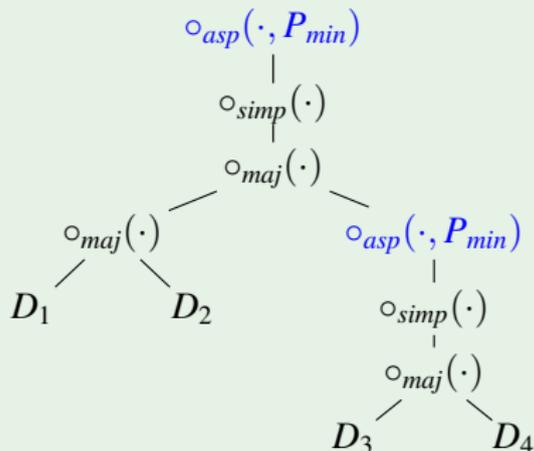
Extended Encoding \hat{E} :

Multiple diagrams within one set of facts: $leaf(L, C) \Rightarrow leaf_{in}(I, L, C)$

- Evaluate P' under ASP semantics

Reasoning about Decision Diagrams

Example: Node Count Minimization



$$\begin{aligned}
 P_{min} = \{ & cnt(I, C) \leftarrow LC = \#count\{L : leaf_{in}(I, L, C)\}, \\
 & IC = \#count\{N : inner_{in}(I, N)\}, \\
 & root_{in}(I, R), C = LC + IC \\
 & c(I) \leftarrow root_{in}(I, R), not \neg c(I) \\
 \neg c(I) \vee \neg c(J) \leftarrow & root_{in}(I, R), root_{in}(J, S), I \neq J \\
 leaf(L, C) \leftarrow & c(I), leaf_{in}(I, L, C) \\
 & \dots \\
 else(N_1, N_2) \leftarrow & c(I), else_{in}(I, N_1, N_2) \\
 \perp \leftarrow M = \#min\{ & NC : cnt(I, NC)\}, \\
 c(I), cnt(I, C), C > M & \}
 \end{aligned}$$

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DNA Classification

Motivation

- Given: **Sequence** over $\{A, C, G, T\}$
- Question: Is it **coding or junk** DNA?

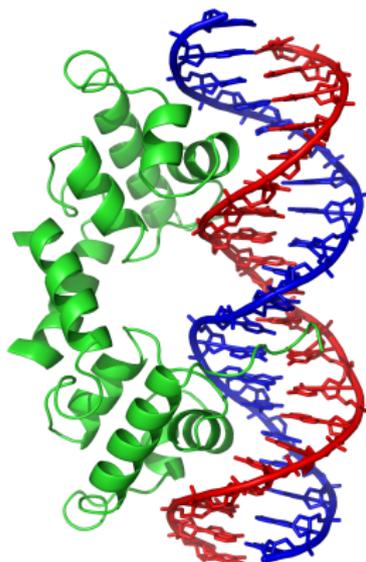
Usual Approach

Training

- 1 Annotated training set
- 2 Compute **statistical features**
- 3 Machine-learning algorithms

Classification

- 1 Compute the same features
- 2 Apply decision diagram



DNA Classification

Advanced Approach [Salzberg et al., 1998]

- Train **multiple** diagrams
varying training sets, algorithms, features, etc.
- Merge them afterwards

Benefits

- Parallelization
- Increase accuracy (cf. genetic algorithms)
- Smaller training set suffices

Hardcoded implementation: **MORGAN system**

DNA Classification

MORGAN's strategy in MELD

- MORGAN's strategy plugged into MELD as **merging operator** \circ_M
- Benefits identified in [5] confirmed

MORGAN vs. MELD-based system

- Not **hardcoded** but **modular**
- Clear separation: **merging operation** / other system components
- **reuse** / **exchange** of the merging operator
- Experiment with **different** merging strategies
- Produce **multiple** diagrams and **reason** about them

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Summary

- MELD: Integration of multiple collections of belief sets
- **Instantiation** for **decision diagram merging**:
 - 1 **Encoding** of decision diagrams as belief sets
 - 2 Special **merging operators** for decision diagrams

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Advantages

- Reuse of operators
- Evaluate different operators empirically
- Automatic recomputation of result
- Release user from routine tasks

Download

URL: <http://www.kr.tuwien.ac.at/research/dlvhex/ddm.html>

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